

Errata

in the book

Klaus Weihrauch, Computable Analysis

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(edited by the author)

Last update: 15 Jan 2005 , correct text is included in $\llbracket \quad \rrbracket$.

- p.VII 1.-11** : $\llbracket \dots$ by A. Turing [Tur36], $\dots \rrbracket$
- p.2 1.-13** : $\llbracket \dots$ R. L. Goodstein [Goo59], $\dots \rrbracket$
- p.2 1.-9** : $\llbracket \dots$ and S. Smale [\dots]
- p.4 1.-4** : $\llbracket \dots (f(n) + 1) \dots \rrbracket$
- p.7 1.-1** : \llbracket **Fig. 1.4.** Characteristic function and distance function of $[0; 1]$. \rrbracket
- p.8 1.9** : replace “boundaries” by \llbracket endpoints \rrbracket .
- p.10 1.17** : $\llbracket x := \sum_{k=n}^{-\infty} a_k \cdot 2^k \rrbracket$
- p.11 1.20** : $\llbracket (i = 1, \dots, k) \rrbracket$
- p.11 1.-18** : replace “ $f[\{a\}] \subseteq \text{dom}(B)$ ” by $\llbracket f[\{a\}] \subseteq \text{dom}(g) \rrbracket$
- p.19 1.-1** : add: \llbracket More general tuplings are defined straightforwardly, e.g., $\langle x, p, z \rangle := \langle \langle x, p \rangle, z \rangle$ for $x, z \in \Sigma^*$ and $p \in \Sigma^\omega$. \rrbracket
- p.20 1.18** : $\llbracket f : \subseteq Y \times Y \times \dots \rrbracket$

- p.30** 1.-10 : insert after “is open”: $\llbracket \text{in } \text{dom}(f) \rrbracket$
- p.30** 1.-8 : $\llbracket \text{computation} \rrbracket$
- p.30** 1.-4 : insert before “contained in”: $\llbracket \text{relative to } \text{dom}(f) \rrbracket$
- p.30** 1.-3 : replace: $\llbracket \text{Therefore, } f^{-1}[w\Sigma^\omega] \text{ is open in } \text{dom}(f). \rrbracket$
- p.45** 1.13 : $\llbracket (\text{union intersection, complement, preimage}). \rrbracket$
- p.46** 1.2 : cut: “and $f^{-1}[W]$ is r.e. open”
- p.52** 1.24 : $\llbracket \text{En}(p) := \{n \in \mathbb{N} \mid 10^{n+1}1 \triangleleft p\}. \rrbracket$
- p.54** 1.9 : $\llbracket , \text{ i.e., } \rrbracket$
- p.54** 1.-7 : $\llbracket \text{Let } \nu \text{ and } \nu' \text{ be notations ... } \rrbracket$
- p.54** 1.-1 : $\llbracket f : \subseteq M \rightarrow M' \rrbracket$
- p.65** 1.5 : $\llbracket \delta_{S'} \equiv \delta_S \rrbracket$
- p.67** 1.11 : replace the proof of 1. by:
 $\llbracket \text{It suffices to prove that } \delta_S^{-1}[A] \text{ is open in } \text{dom}(\delta_S) \text{ for all } A \in \sigma. \text{ We have } \delta_S^{-1}[A] = \{p \in \Sigma^\omega \mid (\exists w \in \nu^{-1}[A]) \iota(w) \triangleleft p\} \cap \text{dom}(\delta_S). \text{ The first set is } \bigcup_{\nu(w)=A} \Sigma^* \iota(w) \Sigma^\omega, \text{ hence an open set. Therefore, } \delta_S^{-1}[A] \text{ is open in } \text{dom}(\delta_S). \rrbracket$
- p.73** 1.13 : replace [Sch00] by $\llbracket \text{[Sch01]} \rrbracket$ which is Schröder, Matthias, Admissible Representations of Limit Spaces, in: Computability and Complexity in Analysis, Blanck, Jens and Bratka, Vasco and Hertling, Peter (eds.), LNCS 2064, 273–295, Springer, Berlin, 2001.
- p.74** 1.11 : $\llbracket \text{Let } \nu \text{ be a notation with r.e. domain. Show that there is some ... } \rrbracket$

- p.86** 1.10 : no linebreak
- p.90** 1.16 : \llbracket the real numbers such that for rational numbers a , the properties ... \rrbracket
- p.91** 1.17 : no linebreak
- p.92** 1.-5 : \llbracket ... infinite base- n fractions ... \rrbracket
- p.97** 1.11 : replace “ (δ, δ) ” by $\llbracket [\delta, \delta] \rrbracket$ in Line 11 (twice)
- p.97** 1.12 : replace the proof as follows:
 \llbracket Assume that the relation “ $x = y$ ” is $[\delta, \delta]$ -open. By Def. 3.1.3.2, $[\delta, \delta]^{-1}\{(x, y) \mid x = y\} = U \cap \text{dom}[\delta, \delta]$ for some open $U \in \Sigma^\omega$. Consider $z = \delta(p)$. Then $\langle p, p \rangle \in U$. Since U is open, $\langle p, p \rangle \in w\Sigma^\omega \subseteq U$ for some $w \in \Sigma^*$. Then $\langle v\Sigma^\omega, v\Sigma^\omega \rangle \subseteq w\Sigma^\omega \subseteq U$ for some $v \in \Sigma^*$ such that $p \in v\Sigma^\omega$. For any $q \in v\Sigma^\omega \cap \text{dom}(\delta)$, $\langle p, q \rangle \in \langle v\Sigma^\omega, v\Sigma^\omega \rangle \subseteq w\Sigma^\omega \subseteq U$, hence $\delta(p) = \delta(q) = x$. Therefore, $\delta[z\Sigma^\omega] = \{z\}$. As we have shown, for every $z \in \mathbb{R}$ there is some $w \in \Sigma^*$ such that $\delta[w\Sigma^\omega] = \{z\}$. But this is impossible, since Σ^* is countable but \mathbb{R} is uncountable. If “ $x \leq y$ ” is $[\delta, \delta]$ -open, then also “ $x \geq y$ ” is $[\delta, \delta]$ -open, hence “ $x = y$ ” is $[\delta, \delta]$ -open by Thm. 2.4.5 (contradiction). Finally, if “ $x < y$ ” is $[\delta, \delta]$ -clopen, then “ $x \geq y$ ” is $[\delta, \delta]$ -open (contradiction). \rrbracket
- p.97** 1.-3 : \llbracket ... substituting the n - \rrbracket
- p.100** 1.-11 : replace [BH00] by \llbracket [BH02] \rrbracket which is Brattka, Vasco and Hertling, Peter, Topological properties of real number representations, TCS 284.2, 2002, 241-257
- p.102** 1.14 : replace 2^n by $\llbracket 2^{n+1} \rrbracket$ (twice)
- p.112** 1.-12 : replace a by $\llbracket c \rrbracket$, also in the figure.
- p.120** 1.-8 : \llbracket ... Definition 3.1.2). \rrbracket

- p.128** 1.-11 : replace “(Exercise 4.1.21)” by [[(Def. 4.1.21 and Exercise 4.1.21).]]
- p.136** 1.-14 : linebreak before [[Since a set ...]]
- p.140** 1.22 : replace ψ by [[$\psi_{<}$]]
- p.142** 1.-5 : [[... is a recursive ...]]
- p.148** insert before 1.11 : [[Notice that κ_H is admissible w.r.t. the Hausdorff topology on \mathcal{K}^* .]]
- p.154** 1.-15 : [[... function is ...]]
- p.157** 1.1 : continue [[for r.e. open sets A . δ_{\rightarrow}^A and δ_{oo}^A are equivalent for r.e. closed sets A .]]
- p.157** 1.-4 : linebreak after [[... translates δ_{oo}^A to δ_{co}^A .]]
- p.172** 1.7 : in the denominators replace $1 + f_i$ by [[$1 + |f_i|$]] and $1 + f_i(x)$ by [[$1 + |f_i(x)|$].]
- p.179** 1.-12, -11 : replace “is non-negative ... $> 1/2$.” by [[is non-negative, computable and has no computable zero. Its set of zeroes has Lebesgue measure $> 1/2$.]]
- p.180** 1.-17 : [[Finding the *position* of the maximum value ...]]
- p.190** 1.14 : [[... its restriction to ...]]
- p.191** 1.17, 19 : [[$\{z \in \mathbb{C} \mid |z| < R\}$]]
- p.209** 1.-15 : [[which on input $p = \iota(u_0)\iota(u_1) \dots \in \text{dom}(\rho^a)$ determines ...]]
- p.228** 1.-3 : replace “every rational” by [[On every compact subset of its domain, every rational]]

- p.232** 1.4 : [... can be encoded by a quadratic boolean matrix ...]
- p.233** 1.-4 : [(see Fig. 7.6).]
- p.238** 1.3 : [First we show that δ is (τ_C, τ) -continuous, where τ is the metric topology on M . Since ...]
- p.238** 1.21 : [... in $\text{dom}(\delta)$ and δ is continuous,]
- p.239** 1.10 : insert before [Definition]: [If $\alpha \equiv \alpha'$ for effective metric spaces $\mathbf{M} = (M, d, A, \alpha)$ and $\mathbf{M}' = (M, d, A, \alpha')$, then $\delta_{\mathbf{M}} \equiv \delta_{\mathbf{M}'}$. Furthermore, for every notation α with r.e. domain there is an equivalent one with recursive domain (Exercise 3.2.15). Therefore, for semi-computable and computable metric spaces we may assume w.l.o.g. that $\text{dom}(\alpha)$ is recursive.]
- p.239** 1.-4 : replace “computable” by [effective]
- p.239** 1.-2 : [... and \mathbf{S} is a computable ...]
- p.239** add to **Thm. 8.1.4** : [3. $\delta_{\mathbf{M}}$ is admissible (w.r.t. the metric topology).]
- p.244** 1.13 : [... the *attractor* of ...]
- p.255** 1.5 : [... the formula in 1. can be ...]
- p.260** 1.-8 : replace “[Str84]” by [[Str87,BCS97]], where [BCS97] is Bürgisser, Peter; Clausen, Michael and Shokrollahi, M. Amin, Algebraic complexity theory, Springer, Berlin, 1997,
- p.261** 1.-4 : replace [Bra99b] by [[Bra03f]] which is Brattka, Vasco, The Emperor’s New Recursiveness: The Epigraph of the Exponential Function in Two Models of Computability, Words, Languages & Combinatorics III, Ito, Masami and Imaoka, Teruo (eds), World Scientific Publishing, Singapore 2003, 63–72.

p.267 1.19 : [[... which corresponds to a given]]

p.277 1.24 : [[$\delta_{\mathbb{B}}$, 52,66,84]]

p.283 1.14 : replace “real RAM” by [[real-RAM]]